

**The Eigenvalue Decomposition**

- Definitions and Basics.
- Localization of Eigenvalues. Sensitivity.
- Applications: Roots of Polynomials. Functions of Matrices.

**Computing Eigenvalues**

- Rayleigh Quotient.
- The Power iteration. Inverse Iteration.

**Remark** Eigenvectors are never unique. If  $x$  is an eigenvector so is  $\alpha x$ . Only the subspaces  $\text{Null}(A - \lambda I)$  are unique.

**Definition** Let  $X = (x_1, x_2, \dots, x_k)$  be all the linearly independent eigenvectors associated with  $A$  and  $D = \text{diag}(\lambda_1, \lambda_2, \dots, \lambda_k)$ . Then  $AX = XD$ .

**Definition** Let  $A \in \mathbb{C}^{n \times n}$ . If there is a scalar  $\lambda \in \mathbb{C}$  and a vector  $x \neq 0$  in  $\mathbb{C}^n$  such that

$$Ax = \lambda x$$

then  $\lambda$  is an *eigenvalue* and  $x$  is an *eigenvector*.

**Remark** Real matrices can have complex eigenvalues. It is easier to treat the complex case.

**Lemma** If  $\lambda$  is an eigenvalue then  $A - \lambda I$  is *singular*. This means that  $\text{Null}(A - \lambda I) \neq \{0\}$ .

**Definition** If  $A \in \mathbb{C}^{n \times n}$  has a full set of linearly independent eigenvectors so  $X = (x_1, \dots, x_n)$ . Then

$$A = XDX^{-1}$$

is called the *eigenvalue decomposition*.

**Remark** In the above case  $A$  is called *non-defective*. In this case the matrix  $X$  provides a *basis* for  $\mathbb{R}^n$ .

**Definition** Let  $x, y \in \mathbb{C}^n$ . The scalar product is

$$(x, y) = x^H y = \sum_{i=1}^n \bar{x}_i y_i,$$

and the *Euclidean norm* is

$$\|x\|_2^2 = x^H x = \sum_{i=1}^n |x_i|^2,$$

where  $x^H$  is the Hermitean transpose of  $x$ .

**Remark** The real case is a special case of the complex one.

**Lemma** If  $A \in \mathbb{C}^{n \times n}$  is *Hermitean*, i.e.  $A = A^H$ , then its eigenvectors are *unitary* so  $X^{-1} = X^H$  and  $A = XDX^H$ .

**Corollary** If  $A$  is real and symmetric, i.e.  $A = A^T$ , then its eigenvectors are *orthogonal* so  $X^{-1} = X^T$  and  $A = XDX^T$ .

**Remark** Both Hermitean and Symmetric matrices have *real* eigenvalues. Anti-Hermitean, i.e.  $A^H = -A$ , have *pure imaginary* eigenvalues.

**Definition** If a matrix  $X \in \mathbb{C}^{n \times n}$  satisfies  $X^H X = I$  then  $X = (x_1, \dots, x_n)$  is *unitary* and its column vectors form an orthonormal basis for  $\mathbb{C}^n$ .

**Lemma** If  $Q$  is *unitary* then  $\|Qx\|_2 = \|x\|_2$ .

**Remark** In the real case the matrix  $Q$  is *orthogonal*.

**Lemma** If  $(\lambda, x)$  is an eigenpair then  $(A - \lambda I)x = 0$ ,  $x \neq 0$ , and therefore the eigenvalues are the roots of

$$p_A(\lambda) = \det(A - \lambda I) = 0,$$

where  $p_A(\lambda)$  is called the *Characteristic polynomial* of  $A$ .

**Remark** If  $A^{-1}$  exists so that  $Ax = b$  has a unique solution for every  $b$ . Then  $A$  is *non-singular*. This is equivalent to zero not being an eigenvalue of  $A$ .

Suppose  $p_A(\lambda)$  is the characteristic polynomial of  $A$ . Since  $p_A(\lambda)$  has  $n$  roots we may write

$$p_A(\lambda) = (\lambda_1 - \lambda)(\lambda_2 - \lambda) \dots (\lambda_n - \lambda).$$

**Definition** The *Algebraic multiplicity*  $\gamma_1(\lambda)$  of  $\lambda$  is its multiplicity as a root of  $p_A(\lambda)$ .

**Definition** The *Geometric multiplicity*  $\gamma_2(\lambda)$  is given by  $\gamma_2(\lambda) = \dim(\text{null}(A - \lambda I))$

**Remark** This is the number of linearly independent eigenvectors associated with  $\lambda$ .

**Lemma** It holds that  $1 \leq \gamma_2(\lambda) \leq \gamma_1(\lambda)$ .

**Definition** If  $\gamma_1(\lambda) = \gamma_2(\lambda)$  for all eigenvalues  $\lambda$  then  $A$  is *non-defective*.

**Remark** In this case we can diagonalize  $A = XDX^{-1}$ . For a *defective eigenvalue* we have  $\gamma_2(\lambda) < \gamma_1(\lambda)$ .

## Example

Consider the *Jordan block*

$$J = \begin{pmatrix} 2 & 1 & 0 \\ 0 & 2 & 1 \\ 0 & 0 & 2 \end{pmatrix}.$$

What are the eigenvalues? Algebraic and Geometric multiplicity?

## Applications

**Example** Suppose  $A$  is non-defective and consider the first order system of ODEs,

$$y'(t) = Ay(t), \quad t > 0, y(0) = b.$$

The solution can be written

$$y(t) = c_1 x_1 e^{\lambda_1 t} + c_2 x_2 e^{\lambda_2 t} + \dots + c_n x_n e^{\lambda_n t}, \quad \text{with, } c = X^{-1}b.$$

**Example** Roots of the polynomial  $p(x) = x^3 + c_2x^2 + c_1x + c_0$  are the eigenvalues of the *companion matrix*

$$\begin{pmatrix} -c_2 & -c_1 & -c_0 \\ 1 & & \\ & 1 & \end{pmatrix}.$$

**Remark** The matlab code `roots` is based on the companion matrix.

The Taylor series representation of the *scalar* function  $f(t)$  is

$$f(t) = \sum_{k=0}^{\infty} \frac{f^{(k)}(0)}{k!} t^k.$$

For any matrix  $A \in \mathbb{R}^{n \times n}$  we define

$$f(A) = \sum_{k=0}^{\infty} \frac{f^{(k)}(0)}{k!} A^k.$$

**Remark** If the power series for  $f(t)$  is absolute convergent for  $|t| < L$  then the series  $f(A)$  is absolute convergent for  $\|A\| < L$ .

**Proposition** If  $A$  is non-defective then

$$f(A) = Xf(D)X^{-1},$$

where  $X$  is the eigenvector matrix,  $D = \text{diag}(\lambda_i)$  are the eigenvalues, and  $f(D) = \text{diag}(f(\lambda_i))$ .

**Remarks** The eigenvalue decomposition offers a cheap and stable way to compute  $f(A)$ .

Matlab has functions `expm`, `cosm`, etc. There is also a function `funm`.

## Localization of Eigenvalues

**Theorem (Gershgorin I)** The eigenvalues of  $A$  are located in the union of the  $n$  discs.

$$|\lambda - a_{ii}| \leq r_i = \sum_{j \neq i} |a_{ij}|, \quad i = 1, 2, \dots, n.$$

**Remark** Since  $\lambda(A)$  and  $\lambda(A^T)$  we can replace  $r_i$  by  $c_i = \sum_{j \neq i} |a_{ji}|$ .

**Theorem (Gershgorin II)** Every isolated subset of discs contains exactly as many eigenvalues as the number of discs.

**Example** Locate the eigenvalues of the matrix

$$A = \begin{pmatrix} 3.3 & 0.4 & -0.7 \\ 0.3 & 2.8 & 0.2 \\ 0.2 & -0.3 & -4 \end{pmatrix}$$

as accurately as possible. Can you conclude that the eigenvalues are real?

## The power method

Suppose  $A$  is real, non-defective,  $|\lambda_1| > |\lambda_2| \geq \dots \geq |\lambda_n|$ , and  $\{x_i\}$  are the eigenvectors.

**Algorithm** Take  $q^{(0)}$  such that  $\|q^{(0)}\|_2 = 1$  and form for  $k = 1, 2, \dots$ ,

$$\begin{aligned} w^{(k)} &= Aq^{(k-1)}, \\ \rho_{k-1} &= (q^{(k-1)})^T w^{(k)}, \\ q^{(k)} &= w^{(k)} / \|w^{(k)}\|_2. \end{aligned}$$

Then  $(\rho_k, q^{(k)})$  converge to the eigenpair  $(\lambda_1, x_1)$ .

**Stopping rule** If  $A = A^T$  and  $r = Aq^{(k)} - \rho_k q^{(k)}$ . Then  $\|r\|_2 < \varepsilon$  ensures that  $|\lambda_1 - \rho_k| < \varepsilon$ .

**Definition** Let  $A \in \mathbb{R}^{n \times n}$  and  $u \in \mathbb{R}^n$  be a non-zero vector. The function

$$\rho(u) = \frac{u^T A u}{u^T u},$$

is called the *Rayleigh quotient*.

**Remark** The Rayleigh quotient is obtained by treating  $Au = \rho u$  as a least squares problem. Thus if  $(x, \lambda)$  is an eigenpair of  $A$  then  $\rho(x) = \lambda$ .

**Proposition** The power method computes estimates  $(\rho_k, q^{(k)})$  of  $(\lambda_1, x_1)$  that satisfy

$$q^{(k)} = \pm x_1 + O(\gamma^k), \quad \text{and,} \quad \rho_k = \lambda_1 + O(\gamma^{k_1}),$$

where  $\gamma = |\lambda_2|/|\lambda_1|$  and  $k_1 = 2k$  if  $A$  is symmetric and  $k_1 = k$  otherwise.

**Remark** The speed of convergence depend on the *quotient*  $\gamma = |\lambda_2/\lambda_1|$ . The factor  $\gamma$  can be improved by linear transformations.

**Lemma** Suppose  $B = A - sI$ . Then  $\lambda(B) = \lambda(A) - s$ .

### Example Suppose

```
>> A=[ 3 4 1 ; 4 5 -1 ; 1 -1 6]
```

The eigenvalues and eigenvectors are

```
>> [X,D]=eig(A)
X =
-0.7674    0.2242   -0.6007   -0.4297         0         0
 0.6046   -0.0587   -0.7943         0    6.2909         0
 0.2134    0.9728    0.0905         0         0    8.1388
D =
```

## Inverse iteration

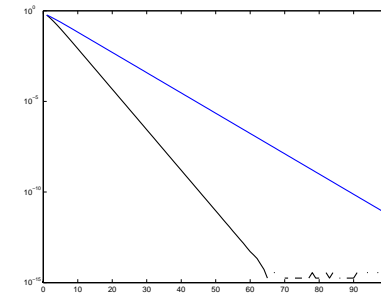
**Proposition** Let  $(\lambda, x)$  be an eigenpair of  $A$ . Put  $B = (A - sI)^{-1}$ . Then  $(\mu, x)$ , with  $\mu = 1/(\lambda - s)$ , is an eigenpair of  $B$ .

**Example** If  $A$  has eigenvalues  $\lambda_1 = -0.4297$ ,  $\lambda_2 = 6.2909$ , and  $\lambda_3 = 8.1388$ . Then with  $s = 8$  we get,

$$\gamma = \left| \frac{\lambda_2 - s}{\lambda_3 - s} \right| = \left| \frac{6.2909 - 8}{8.1388 - 8} \right| \approx 0.0812$$

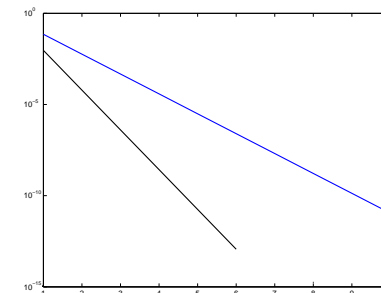
Much faster convergence than the power method ( $\gamma = |\lambda_2/\lambda_3| = 0.7729$ ).

### Example Power iteration results.



The errors  $\|\bar{x}^{(k)} - x_3\|_2$  (blue) and  $|\rho_k - \lambda_3|$  (black).  $A$  is symmetric and  $\gamma = |\lambda_2/\lambda_3| = 0.7729$ . This is *slow* convergence.

### Example Inverse iteration results with $s = 8$ .



The errors  $\|\bar{x}^{(k)} - x_3\|_2$  (blue) and  $|\rho_k - \lambda_3|$  (black). This is *fast* convergence.

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**Lemma** Let  $(x, \lambda)$  be an eigenpair of a symmetric matrix  $A$ , and  $\bar{x}$  an approximation of  $x$ . If  $\|x - \bar{x}\|_2 = O(\varepsilon)$  then  $|\lambda - \rho(\bar{x})| = O(\varepsilon^2)$ .

**Lemma** If  $A$  is symmetric then  $B = (A - sI)^{-1}$  is also symmetric.

**Remark** This means the faster convergence.

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**Lemma** If  $\lambda$  is an eigenvalue then  $A - \lambda I$  is *singular*. We can compute the eigenvector by, e.g., setting  $x_1 = 1$ , and solving  $Ax = \lambda x$ .

**Remark** This means that we only need to compute eigenvalues. Eigenvectors can easily be obtained.

**Question** How to find more efficient methods for computing eigenvalues.