

Lecture 1

①
P.6

$$C_{ij} = \sum_{k=1}^n a_{ik} b_{kj}$$

Load $B(:,j), C(:,j)$

for $k=1:n$

load $A(i,k)$

$$C(i,j) = C(i,j) + A(i,k) \times B(k,j)$$

Total $O(n^3)$ multiplications/memory access

②

Change from column to block
P.8 storage requires $n \times p = n\sqrt{n} = n^{1.5}$ //
memory load/store

The algorithm is

Load $C_{ij} \in \mathbb{R}^{p \times p}$

for $k=1:p$

load $A_{ik}, B_{kj} \in \mathbb{R}^{p \times p}$

$$C_{ij} = C_{ij} + A_{ik} \cdot B_{kj}$$

$O(p^3)$ mults.

Block size is $p = \sqrt{n}$ so work

$$p^2 \cdot p \cdot p^3 = p^6 = (\sqrt{n})^6 = n^3 \text{ mult.} //$$

and need

$$p^2 \cdot p \cdot 2 = 2p^3 = 2n^{1.5} \text{ memory loads}$$

③

P.12

$$y = \underbrace{(AB)}_{\substack{n^3 \\ \text{or } n^2}} x = A \underbrace{(Bx)}_{n^2}$$

Either $n^3 + n^2$ or $2n^2$ multiplications.

Matlab evaluates from left to right

>> $n = 10^3$; $A = \text{rand}(n, n)$; $B = \text{rand}(n, n)$;

>> $x = \text{rand}(n, 1)$

>> tic, $y = A * B * x$; toc $\Rightarrow 4.4$ s

>> tic, $y = A * (B * x)$; toc $\Rightarrow 0.002$ s

④

P.13

$$A_{ij} = u_i \cdot v_j \quad 1 \leq i, j \leq n \Rightarrow n^2 \text{ multiply. scalar.}$$

Example

$$y = (uv^T)x = u \cdot \underbrace{(v^T x)}_{\substack{n \text{ op}}} \\ \underbrace{}_{n \text{ op}}$$

Need $O(n)$ multiplications and n op memory.

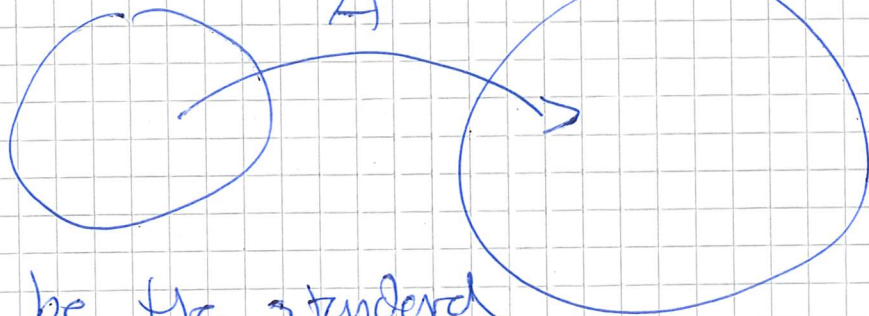
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P.16

$\mathbb{R}^n$

A

$\mathbb{R}^m$



Let $\{e_i\}$ be the standard

basis and $A = (a_1, a_2, \dots, a_n)$, $a_k \in \mathbb{R}^m$

Then

$Ae_i = a_i \in \text{Range}(A)$ for all i .

So

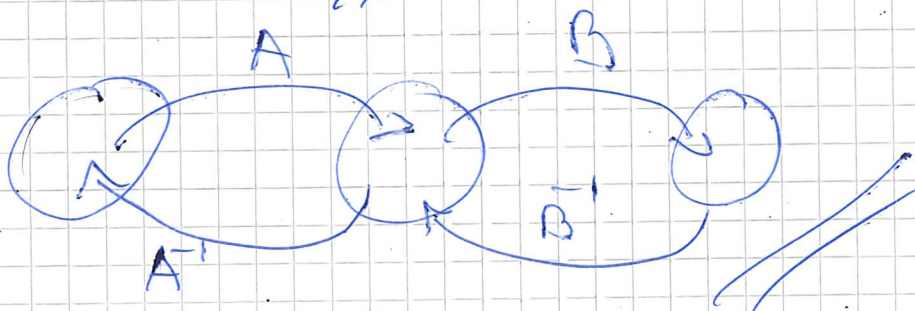
$\text{Range}(A) = \text{span}(a_1, \dots, a_n) \subset \mathbb{R}^m$

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P.17

$$\begin{aligned} I &= (AB)^{-1}(AB) = B^{-1}A^{-1}AB = B^{-1}(A^{-1}A)B \\ &= B^{-1}B = I // \end{aligned}$$

or



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P.18

$A \in \mathbb{R}^{m \times n}$. If $b \in \text{Range}(A)$ there is an $x \in \mathbb{R}^n$
so $Ax = b$. This is existence.

If $\dim(\text{Range}(A)) = m$ then $b \in \mathbb{R}^m \equiv \text{Range}(A)$

so existence for all $b \in \mathbb{R}^m$ //

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P.18

If $Ax=b$ has a solution x_1 and $x_2 \in \text{Null}(A)$

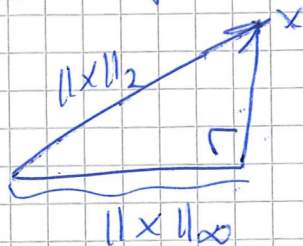
$$\text{then } A(x_1 + x_2) = \underbrace{Ax_1}_{=b} + \underbrace{Ax_2}_{=0} = b$$

The solution is unique if $\text{null}(A) = \{0\}$

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P.19

Example $\|x\|_1 = \sum_{i=1}^n |x_i|$



$$\|x\|_\infty \leq \|x\|_2 \leq \|x\|_1$$

Many different norms and relations //

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P.20

$$\|A\|_\infty = \max_{1 \leq i \leq n} \sqrt{\lambda_i(A^T A)}$$

Demerding to compute

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P.22

$$\|AB\| \leq \max_{x \neq 0} \frac{\|ABx\|}{\|x\|} = \max_{x \neq 0} \frac{\|A(Bx)\|}{\|Bx\|} \frac{\|Bx\|}{\|x\|}$$

$$\leq \|B\| \max_{y \neq 0} \frac{\|Ay\|}{\|y\|} = \|B\| \cdot \|A\| //$$