

Lecture 2

- ① proof Have to show that $Ax=b$ has a unique solution for all $b \in \mathbb{R}^n$. If x_1 and x_2 are two solutions then $A(x_1 - x_2) = Ax_1 - Ax_2 = b - b = 0$ so $x_1 - x_2 \in \text{null}(A) = \{0\} \Rightarrow x_1 = x_2$ so uniqueness.
- Also: $\text{range}(A) = \mathbb{R}^n$ due to previous lemma.

② $A = \begin{pmatrix} 10^{-6} & & \\ & 10^{-6} & \\ & & 10^{-6} \end{pmatrix} \Rightarrow \det(A) = 10^{-18} < \mu =$ machine precision.

③ Row operation $(1 \ 2 \ 2 \mid 3) - 0.25(4 \ 4 \ 2 \mid 6)$
 $= (0 \ 1 \ 1.5 \mid 1.5) //$

④ $M^{-1} = \begin{pmatrix} 1 & 0 & 0 \\ m_{21} & 1 & 0 \\ m_{31} & 0 & 1 \end{pmatrix}$ Add back the same multiple of the first row.

⑤ $P_{23} M_1^{-1} P_{23} = P_{23} \begin{pmatrix} 1 & 0 & 0 \\ m_{21} & 1 & 0 \\ m_{31} & 0 & 1 \end{pmatrix} P_{23}^T = P_{23} \begin{pmatrix} 1 & 0 & 0 \\ m_{21} & 0 & 1 \\ m_{31} & 1 & 0 \end{pmatrix} \begin{matrix} \uparrow \\ \downarrow \end{matrix}$
 $= \begin{pmatrix} 1 & 0 & 0 \\ m_{31} & 1 & 0 \\ m_{21} & 0 & 1 \end{pmatrix} \hat{M}_1^{-1} //$

also $\hat{M}_1^{-1} \hat{M}_2^{-1} = \hat{M}_1^{-1} \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & m_{32} & 1 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 0 \\ m_{31} & 1 & 0 \\ m_{21} & m_{32} & 1 \end{pmatrix} //$

(6)

If $PA=LU$ then

P.10

$$\det(A) = \det(P^{-1}) \det(L) \det(U) =$$

$$= (\pm 1) \cdot 1 \cdot \prod_{i=1}^n u_{ii} //$$

If Gaussian elimination works then A is non singular.

(7)

$$PAx = Pb \Rightarrow LUx = Pb \Rightarrow$$

P.10

$$x = U^{-1}L^{-1}(Pb) = "U \setminus (L \setminus (P * b))";$$

(8)

No pivoting means

$$A = LU = L(D\tilde{L}^T) = (LD^{1/2})(\tilde{L}D^{1/2})^T$$

$$= R^T \tilde{R}$$

$$A \text{ symmetric} \Rightarrow R = \tilde{R} //$$

(9)

proof

$$\text{Have } A(x + \delta x) = b + \delta b \Rightarrow \delta x = \bar{A}^{-1} \delta b$$

$$\text{and } \|\delta x\| \leq \|\bar{A}^{-1}\| \cdot \|\delta b\|. \text{ Also } \|b\| \leq \|A\| \cdot \|x\|$$

Thus

$$\frac{\|\delta x\|}{\|x\|} \leq \frac{\|\bar{A}^{-1}\| \cdot \|\delta b\|}{\|b\| / \|A\|} = \underbrace{\|A\| \|\bar{A}^{-1}\|}_{K(A)} \cdot \frac{\|\delta b\|}{\|b\|} //$$

Valid for $\|\cdot\|_2$, $\|\cdot\|_1$ and $\|\cdot\|_\infty$ //

(10) In Matlab $\mu \approx 1.1 \cdot 10^{-16}$ is the machine precision.

P.13 Best case. $\|r\| \approx 10^{-14} - 10^{-16}$

$$\kappa_2(A) \approx 10^4 - 10^{12}$$

$$\text{Relative error } 10^{-12} - 10^{-2}$$

Rough estimate of the error.

(11) $\|Ax - b\|_2$ is minimized if $r \perp \text{Range}(A)$

P.14

$$\text{or } (b - Ax)^T A = 0 \Leftrightarrow \{A^T A x = A^T b\}$$

The proof is based on having a scalar product

(12) Modell: $p(t_k) = c_0 + c_1 t_k + c_2 t_k^2, 1 \leq k \leq m$

P.17 or

$$\begin{pmatrix} 1 & t_1 & t_1^2 \\ 1 & t_2 & t_2^2 \\ \vdots & \vdots & \vdots \\ 1 & t_m & t_m^2 \end{pmatrix} \begin{pmatrix} c_0 \\ c_1 \\ c_2 \end{pmatrix} = \begin{pmatrix} y_1 \\ y_2 \\ \vdots \\ y_m \end{pmatrix} \quad \{m \times 3\}$$

(13) let $Q = (q_1, q_2, \dots, q_m)$ then for $x \in \mathbb{R}^m$

$$y = Q Q^T x = Q \begin{pmatrix} q_1^T x \\ \vdots \\ q_m^T x \end{pmatrix} = (q_1^T x) q_1 + \dots + (q_m^T x) q_m$$

Orthogonal projection on $\text{span}(q_1, \dots, q_n)$ //

(14)

P. 21

$$\begin{array}{ccccccc}
 \boxed{A} & = & \boxed{Q} & \boxed{R} & = & \boxed{Q_1} & \boxed{R} \\
 m \times n & & m \times m & m \times n & & m \times n & n \times n
 \end{array}$$

To store Q requires m^2 memory slots. //

$$(15) \quad \|Ax - b\|_2^2 = \left\| \begin{pmatrix} R \\ 0 \end{pmatrix} x - b \right\|_2^2 = \left\| \begin{pmatrix} Rx - Q_1^T b \\ 0 - Q_2^T b \end{pmatrix} \right\|_2^2$$

$Q = (Q_1 \ Q_2)$

$$= \|Rx - Q_1^T b\|_2^2 + \|Q_2^T b\|_2^2$$

basis for
Range(A)

basis for
null(A)

so minimum for $x = R^{-1} Q_1^T b //$