

Lecture 4

① proof (x, λ) eigenpair $\Rightarrow Ax = \lambda x \Rightarrow (A - \lambda I)x = 0$
 $x \neq 0$ so non-trivial null space.

② $AX = A(x_1, \dots, x_n) = (Ax_1, \dots, Ax_n) = (\lambda_1 x_1, \dots, \lambda_n x_n)$
 $= (x_1, \dots, x_n) \begin{pmatrix} \lambda_1 & & \\ & \ddots & \\ & & \lambda_n \end{pmatrix}$

$$\begin{matrix} \boxed{n \times n} & \boxed{n \times k} \\ A & X \end{matrix} = \begin{matrix} \boxed{n \times k} & \boxed{k \times k} \\ X & D \end{matrix}$$

$$\text{If } X \in \mathbb{R}^{n \times n} \Rightarrow AX = XD \Rightarrow A = XD X^{-1}$$

$$\textcircled{3} (A^H)_{ij} = \overline{A_{ji}} //$$

④ proof If $A^H = -A$ and $Ax = \lambda x$ then

$$\begin{aligned} \lambda &= \lambda x^H x = x^H (\lambda x) = x^H A x = x^H (-A^H x) = -(Ax)^H x \\ &= -(\lambda x)^H x = -\overline{\lambda} x^H x = -\overline{\lambda} \end{aligned}$$

$$\text{So } \lambda = -\overline{\lambda} \text{ and } \text{Re}\{\lambda\} = 0 //$$

$$\textcircled{5} P_3(\lambda) = \det(3 - \lambda I) = (2 - \lambda)^3 \Rightarrow \lambda = 2 //$$

We have $\delta_1(2) = 3$. Also $x = \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix}$ is the only eigenvector $\Rightarrow \delta_2(2) = 1$.

Jordan decomposition $A = Q D Q^T$, $D = \text{block diagonal}$.

⑥ $p(\lambda) = \det(A - \lambda I) = \det \begin{pmatrix} -c_2 - \lambda & -c_1 & -c_0 \\ 1 & -\lambda & 0 \\ 0 & 1 & -\lambda \end{pmatrix}$

$$= -\det \begin{pmatrix} -c_2 - \lambda & -c_1 \\ 1 & -\lambda \end{pmatrix} (-\lambda) - \det \begin{pmatrix} 1 & -\lambda \\ 0 & 1 \end{pmatrix} (-c_0)$$

$$= c_0 + \lambda \det \begin{pmatrix} -c_2 - \lambda & -c_1 \\ 1 & -\lambda \end{pmatrix} = \dots = \lambda^3 + c_2 \lambda^2 + c_1 \lambda + c_0$$

Rekursion step
in general proof.

⑦ proof $f(A) = \sum_{k=0}^{\infty} c_k A^k = \left\{ \begin{array}{l} A = XDX^{-1} \Rightarrow \\ A^k = XDX^{-1}XD^{-1}XD^{-1} \dots XDX^{-1} \end{array} \right\}$

$$= \sum_{k=0}^{\infty} c_k X D^k X^{-1} = X \left(\sum_{k=0}^{\infty} c_k D^k \right) X^{-1} =$$

$$= X \begin{pmatrix} \sum_{k=0}^{\infty} c_k \lambda_1^k & & \\ & \ddots & \\ & & \sum_{k=0}^{\infty} c_k \lambda_n^k \end{pmatrix} X^{-1} = X f(D) X^{-1}$$

$\underbrace{\sum_{k=0}^{\infty} c_k \lambda_n^k}_{f(\lambda_n)}$

⑧ proof Let (λ, x) be an eigenpair of A . Then

$$Ax = \lambda x \Leftrightarrow (\lambda - a_{ii})x_i = \sum_{j \neq i} a_{ij}x_j, \quad i=1, 2, \dots, n$$

If $|x_n| = \max |x_i|$ then

$$|\lambda - a_{nn}| \leq \sum_{j \neq n} |a_{nj}| \left| \frac{x_j}{x_n} \right| \leq \sum_{j \neq n} |a_{nj}| = r_n //$$

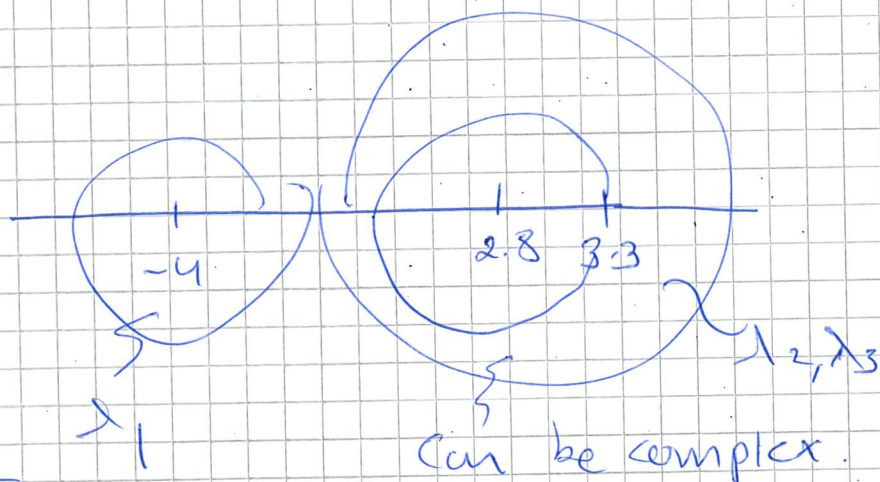
So λ belongs to one disc //

⑨ DBCS

$$|3.3 - \lambda| \leq 1.1$$

$$|2.8 - \lambda| \leq 0.5$$

$$|-4 - \lambda| \leq 0.5$$



looking at A^T gives same results

⑩ If $A \bar{x} \approx \lambda \bar{x} \Rightarrow \underbrace{\bar{x}}_{n \times 1} \cdot \underbrace{\lambda}_{1 \times 1} = \underbrace{A \bar{x}}_{n \times 1}$

Normal equations: $\bar{x}^T \bar{x} \cdot \lambda = \bar{x}^T A \bar{x} \Rightarrow \lambda = \frac{\bar{x}^T A \bar{x}}{\bar{x}^T \bar{x}}$

⑪ let $A = XDX^T$, $D = \text{diag}(\lambda_i)$, $|\lambda_1| > |\lambda_2| \geq \dots \geq |\lambda_n|$

X is a basis for $\mathbb{R}^n$ so any vector can be written

$$x = Xc = \sum_{i=1}^n c_i x_i, \quad X = (x_1, x_2, \dots, x_n)$$

so $A^k x = \sum_{i=1}^n c_i A^k x_i = \sum_{i=1}^n c_i \lambda_i^k x_i = \lambda_1^k \left(c_1 x_1 + \sum_{i=2}^n c_i \left(\frac{\lambda_i}{\lambda_1} \right)^k x_i \right)$

$$\Rightarrow \lambda_1^k c_1 x_1 //$$

Power iteration

$$x^{(k)} = A^k x; \quad q^{(k)} = \frac{x^{(k)}}{\|x^{(k)}\|_2}$$

$$\Rightarrow q^{(k)} \rightarrow \pm x_1, \text{ and } e_k = (q^{(k)})^T A q^{(k)} \rightarrow \lambda_1$$

Error is dominated by $c_2 \left(\frac{\lambda_1}{\lambda_2} \right)^k x_2 //$

(12) proof If $Ax = \lambda x$ then

$$(A - sI)x = (\lambda - s)x \Leftrightarrow \underbrace{\frac{1}{\lambda - s}}_{\mu} x = \underbrace{(A - sI)^{-1}}_B x$$

Apply power method to B instead.

Convergence rate

$$\frac{\mu_2}{\mu_1} = \frac{\frac{1}{(\lambda_i - s)}}{\frac{1}{(\lambda_j - s)}} = \frac{\lambda_j - s}{\lambda_i - s}$$

Very fast if $s \approx \lambda_i$ //
